

# Proof of a Conjecture on the Sequence of Exceptional Numbers, Classifying Cyclic Codes and APN Functions

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## Abstract

We prove a conjecture that classifies exceptional numbers. This conjecture arises in two different ways, from cryptography and from coding theory. An odd integer  $t \geq 3$  is said to be exceptional if  $f(x) = x^t$  is APN (Almost Perfect Nonlinear) over  $\mathbb{F}_{2^n}$  for infinitely many values of  $n$ . Equivalently,  $t$  is exceptional if the binary cyclic code of length  $2^n - 1$  with two zeros  $\omega, \omega^t$  has minimum distance 5 for infinitely many values of  $n$ . The conjecture we prove states that every exceptional number has the form  $2^i + 1$  or  $4^i - 2^i + 1$ .

*Key words:* Absolutely irreducible polynomial, coding theory, cryptography.

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## 1 Introduction

The sequence of numbers of the form  $2^i + 1$  or  $4^i - 2^i + 1$  (where  $i \geq 1$ ) is

$$3, 5, 9, 13, 17, 33, 57, 65, 129, 241, 257, 513, 993, 1025, \dots$$

This is sequence number A064386 in the On-Line Encyclopedia of Integer Sequences. It has been known for almost 40 years that these numbers are *exceptional* numbers, in the sense we will define shortly. No further exceptional numbers were found, and it was conjectured that this sequence is the complete list of exceptional numbers. In this article we complete the proof of this conjecture. Somewhat surprisingly, the sequence of exceptional numbers arises in two different contexts, as explained in the excellent survey article of Dillon [2]. We now proceed to give these two different motivations for the conjecture.

### 1.1 Coding theory

We fix our base field  $\mathbb{F}_2$ . Let  $w$  be a primitive  $(2^n - 1)$ -th root of unity in an extension of  $\mathbb{F}_2$ , i.e., a primitive element of  $\mathbb{F}_{2^n}$ . For every odd  $t \geq 3$ , we define  $C_n^t$  as the cyclic code over  $\mathbb{F}_2$  of length  $2^n - 1$  with two zeros  $w, w^t$ . It is well known that if  $t = 3$ , the code  $C_n^3$  has minimum distance 5 for every  $n \geq 3$ . This code is called the 2-error-correcting BCH code. We want to find other values of  $t$  (fixed with respect to  $n$ ) for which the code  $C_n^t$  has minimum distance 5 for infinitely many values of  $n$ . Those values of  $t$  having this property are called **exceptional**. The only known exceptional values for  $t$  are numbers of the form  $t = 2^i + 1$  (known in the field of coding theory as Gold numbers) and  $t = 4^i - 2^i + 1$  (known as Kasami-Welch numbers). We give more on the precise history in Section 2.1. The conjecture stated by Janwa-McGuire-Wilson [5] is

**Conjecture 1:** *The only exceptional values for  $t$  are the Gold and Kasami-Welch numbers.*

Equivalently, the conjecture says that for a fixed odd  $t \geq 3$ ,  $t \neq 2^i + 1$  or  $t \neq 4^i - 2^i + 1$ , the codes  $C_n^t$  of length  $2^n - 1$  have codewords of weight 4 for all but for finitely many values of  $n$ . In this paper we prove Conjecture 1.

## 1.2 Cryptography

The second approach to this problem comes from cryptography. One of the desired properties for an S-box used in a block cipher is to have the best possible resistance against differential attacks, i.e., any given plaintext difference  $a = y - x$  provides a ciphertext difference  $f(y) - f(x) = b$  with small probability. More formally, a function  $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$  is said to be APN (Almost Perfect Nonlinear) if for every  $a, b \in \mathbb{F}_{2^n}$  with  $a \neq 0$  we have

$$\#\{x \in \mathbb{F}_{2^n} \mid f(x+a) + f(x) = b\} \leq 2.$$

Over a field of characteristic 2, APN functions provide optimal resistance to differential cryptanalysis.

Monomial functions  $f(x) = x^t$  from  $\mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$  are often considered for use in applications. The exponent  $t$  is called **exceptional** if  $f(x) = x^t$  is APN on infinitely many extension fields of  $\mathbb{F}_2$ . The conjecture stated by Dillon [2] is

**Conjecture 2:** *The only exceptional exponents are the Gold and Kasami-Welch numbers.*

Of course, Dillon knew that Conjecture 2 is the same as Conjecture 1, as we explain below. Conjecture 2 says that for a fixed odd  $t \geq 3$ ,  $t \neq 2^i + 1$  or  $t \neq 4^i - 2^i + 1$ , the function  $f(x) = x^t$  is APN on at most a finite number of fields  $\mathbb{F}_{2^n}$ . In this paper we prove Conjecture 2.

## 1.3 Summary of Paper

In Section 2 we will explain why Conjecture 1 and Conjecture 2 are the same. Section 3 gives some known results and some background theory that we will need. The proof naturally splits into two cases. We use the notation  $t = 2^i\ell + 1$ , where  $i \geq 1$ , and  $\ell \geq 3$  is odd. The two cases are dependent on the value of  $\gcd(\ell, 2^i - 1)$ . In Section 4 we recall a result of Jedlicka, which proves the result in the case  $\gcd(\ell, 2^i - 1) < \ell$ . In Section 5 we give a proof of the main theorem of this paper, Theorem 16, which proves the case  $\gcd(\ell, 2^i - 1) = \ell$ . In Section 6 we give a counterexample to Conjecture 3 (stated below).

## 2 Background

This proofs in this paper concern the absolute irreducibility of certain polynomials. In this section we will outline how these polynomials arise from Conjectures 1 and 2.

### 2.1 Coding Theory

It is well known that codewords of weight 4 in  $C_n^t$  are equivalent to the polynomial

$$f_t(x, y, z) = x^t + y^t + z^t + (x + y + z)^t \quad (1)$$

having a rational point  $(\alpha, \beta, \gamma)$  over  $\mathbb{F}_{2^n}$  with distinct coordinates. Notice that  $x + y$ ,  $x + z$  and  $y + z$  divide  $f_t(x, y, z)$ , so we may restrict ourselves to rational points of the homogeneous polynomial

$$g_t(x, y, z) = \frac{f_t(x, y, z)}{(x + y)(x + z)(y + z)}. \quad (2)$$

Janwa-Wilson [6] provide the following result using the Weil bound.

**Proposition 1** *If  $g_t(x, y, z)$  has an absolutely irreducible factor defined over  $\mathbb{F}_2$  then  $g_t(x, y, z)$  has rational points  $(\alpha, \beta, \gamma) \in (\mathbb{F}_{2^n})^3$  with distinct coordinates for all  $n$  sufficiently large.*

The following conjecture was proposed by Janwa-McGuire-Wilson.

**Conjecture 3:** *The polynomial  $g_t(x, y, z)$  is absolutely irreducible for all  $t$  not of the form  $2^i + 1$  or  $4^i - 2^i + 1$ .*

We give a counterexample (found with MAGMA) to Conjecture 3 in Section 6. A slightly weaker form of Conjecture 3 is:

**Conjecture 3':** *The polynomial  $g_t(x, y, z)$  has an absolutely irreducible factor defined over  $\mathbb{F}_2$  for all  $t$  not of the form  $2^i + 1$  or  $4^i - 2^i + 1$ .*

By Proposition 1 and the discussion above, it is clear that Conjecture 3  $\Rightarrow$  Conjecture 3'  $\Rightarrow$  Conjecture 1. In this paper we will prove Conjecture 3', and as a result, we prove Conjecture 1.

Notice that  $g_t(x, y, z)$  has no singular points at the infinity, thus the latter conjecture may be reformulated using  $g_t(x, y, 1)$  instead of  $g_t(x, y, z)$ . We write  $f_t(x, y)$  for  $f_t(x, y, 1)$ , and we write  $g_t(x, y)$  for  $g_t(x, y, 1)$ .

For the known exceptional values of  $t$ , that is, when  $t$  has the form  $2^i + 1$  or  $4^i - 2^i + 1$ , the polynomial  $g_t(x, y)$  is known to *not* be absolutely irreducible, and the factorization is described in [6]. We also remark that for some values of  $t$ , such as  $t = 7$ ,  $g_t(x, y)$  is nonsingular and therefore absolutely irreducible, but it is false that  $g_t(x, y)$  is nonsingular for all  $t$  not of the form  $2^i + 1$  or  $4^i - 2^i + 1$ .

## 2.2 Cryptography

It is well known that

$$h_t(x, y) = \frac{(x+1)^t + x^t + (y+1)^t + y^t}{(x+y)(x+y+1)}.$$

has no rational points over  $\mathbb{F}_{2^n}$  besides those with  $x = y$  and  $x = y + 1$  if and only if  $x^t$  is APN over  $\mathbb{F}_{2^n}$ .

Analogous to Proposition 1, Jedlicka [7] showed that as a consequence of the Weil bound we have the following result.

**Proposition 2** *If  $h_t(x, y)$  has an absolutely irreducible factor over  $\mathbb{F}_2$  then  $h_t(x, y)$  has rational points over  $\mathbb{F}_{2^n}$  besides those with  $x = y$  and  $x = y + 1$  for all  $n$  sufficiently large.*

The following conjectures are essentially stated in [7].

**Conjecture 4:** *The polynomial  $h_t(x, y)$  is absolutely irreducible polynomial for all  $t$  not of the form  $2^i + 1$  or  $4^i - 2^i + 1$ .*

A slightly weaker version of this conjecture is:

**Conjecture 4':** *The polynomial  $h_t(x, y)$  has an absolutely irreducible factor defined over  $\mathbb{F}_2$  for all  $t$  not of the form  $2^i + 1$  or  $4^i - 2^i + 1$ .*

By Proposition 2 and the discussion above, it is clear that Conjecture 4  $\Rightarrow$  Conjecture 4'  $\Rightarrow$  Conjecture 2. In this paper we will prove Conjecture 4', and as a result, we prove Conjecture 2. We give a counterexample to Conjecture 4.

### 2.3 Putting the Problems Together

The following is well known to researchers in the area.

**Lemma 3** *Conjecture 3 is true iff Conjecture 4 is true. Conjecture 3' is true iff Conjecture 4' is true.*

Proof: Factoring out  $y^t$  from  $(x+1)^t + x^t + (y+1)^t + y^t$  and letting  $X = \frac{x+1}{y}$  and  $Y = \frac{x}{y}$  gives

$$(x+1)^t + x^t + (y+1)^t + y^t = y^t[X^t + Y^t + 1 + (X+Y+1)^t].$$

Therefore, we can study the irreducibility of  $h_t(x, y)$  or that of  $g_t(x, y)$ , they are equivalent.  $\square$

We can say even more: the monomial  $x^t$  is an APN function over  $\mathbb{F}_{2^n}$  if and only if the code  $C_n^t$  has minimum distance 5. This shows that Conjecture 1 is true iff Conjecture 2 is true.

Conjecture 3	$\iff$	Conjecture 4
$\Downarrow$		$\Downarrow$
Conjecture 3'	$\iff$	Conjecture 4'
$\Downarrow$		$\Downarrow$
Conjecture 1	$\iff$	Conjecture 2

In this paper we will prove Conjecture 3', see the table below. This is equivalent to proving Conjecture 4', and so implies both Conjectures 1 and 2. However, we give a counterexample in Section 6 that shows that Conjectures 3 and 4 are false in general.

**Notation:** Throughout we will let  $t = 2^i \ell + 1$ , where  $i \geq 1$ , and  $\ell \geq 3$  is odd.

We use the notation  $Sing(g_t)$  to denote the set of all singular points of  $g_t$ .

The following box summarizes known results before this paper, and what is done in this paper.

Function	Exceptional	Constraints	Author
$x^{2^i+1}$	Yes	APN iff $(i, n) = 1$	Gold [4]
$x^{4^i-2^i+1}$	Yes	APN iff $(i, n) = 1$	van Lint-Wilson [10], Janwa-Wilson [6]
$x^t$	No	$t \equiv 3(\text{mod } 4), t > 3$	Janwa-McGuire-Wilson [5]
$x^{2^i\ell+1}$	No	$\gcd(\ell, 2^i - 1) < \ell$	Jedlicka [7]
$x^{2^i\ell+1}$	No	$\gcd(\ell, 2^i - 1) = \ell$	This paper

Janwa-McGuire-Wilson proved the  $i = 1$  case. The full proof of Conjecture 3' divides into two cases, according as  $\gcd(\ell, 2^i - 1) < \ell$  or  $\gcd(\ell, 2^i - 1) = \ell$ . Jedlicka proved the case  $\gcd(\ell, 2^i - 1) < \ell$ . In the present work we give a proof of Conjecture 3' in the remaining case when  $\gcd(\ell, 2^i - 1) = \ell$ . We show that Conjecture 3 is false in general. This completes the classification of exceptional exponents.

### 3 Singularities and Bezout's Theorem

Consider  $P = (\alpha, \beta)$ , a point in the plane. Write

$$f_t(x + \alpha, y + \beta) = F_0 + F_1 + F_2 + F_3 + \cdots$$

where  $F_m$  is homogeneous of degree  $m$ . The multiplicity of  $f_t$  at  $P$  is the smallest  $m$  with  $F_m \neq 0$ , and is denoted by  $m_P(f_t)$ . In this case,  $F_m$  is called the tangent cone.

Recall the notation that  $t = 2^i\ell + 1$ , where  $i \geq 1$ , and  $\ell \geq 3$  is odd.

We let  $\lambda = \alpha + \beta + 1$ , then straightforward calculations [6] give

$$F_0 = \alpha^t + \beta^t + \lambda^t + 1$$

$$F_1 = (\alpha^{t-1} + \lambda^{t-1})x + (\beta^{t-1} + \lambda^{t-1})y$$

$$F_{2^i} = (\alpha^{t-2^i} + \lambda^{t-2^i})x^{2^i} + (\beta^{t-2^i} + \lambda^{t-2^i})y^{2^i}$$

$$F_{2^i+1} = (\alpha^{t-2^i-1} + \lambda^{t-2^i-1})x^{2^i+1} + (\beta^{t-2^i-1} + \lambda^{t-2^i-1})y^{2^i+1} + \lambda^{t-2^i-1}(x^{2^i}y + xy^{2^i})$$

and  $F_j = 0$  for  $1 < j < 2^i$ . A point  $P = (\alpha, \beta)$  is singular if and only if  $F_0 = F_1 = 0$ , which happens if and only if  $\alpha, \beta$  and  $\lambda$  are  $\ell$ -th roots of unity (see [6]). We distinguish three types of singular point.

- (I)  $\alpha = \beta = \lambda = 1$ .
- (II) Either  $\alpha = 1$  and  $\beta \neq 1$ , or  $\beta = 1$  and  $\alpha \neq 1$ , or  $\alpha = \beta \neq 1$  and  $\lambda = 1$ .  
 We divide these singular points into two cases:
  - (II.A) Where II holds and  $\alpha, \beta \in GF(2^i)$
  - (II.B) Where II holds and  $\alpha, \beta$  not both in  $GF(2^i)$ .
- (III)  $\alpha \neq 1, \beta \neq 1$  and  $\alpha \neq \beta$ .  
 We divide these singular points into two cases:
  - (III.A) Where III holds and  $\alpha, \beta \in GF(2^i)$
  - (III.B) Where III holds and  $\alpha, \beta$  not both in  $GF(2^i)$ .

Now we summarize some properties already known, for more details see [5].

**Lemma 4** *If  $F_{2^i} \neq 0$  then  $F_{2^i} = (Ax + By)^{2^i}$  where  $A^{2^i} = \alpha^{1-2^i} + \lambda^{1-2^i}$  and  $B^{2^i} = \beta^{1-2^i} + \lambda^{1-2^i}$ .*

The proof is obvious, because we are in characteristic 2. The importance of this lemma is that there is only one distinct linear factor in  $F_{2^i}$ . Another useful fact is that the opposite is true for  $F_{2^{i+1}}$ , as shown in [5]:

**Lemma 5**  *$F_{2^{i+1}}$  has  $2^i + 1$  distinct linear factors.*

### 3.1 Classification of Singularities

The next step is to describe how many singularities of each type there are, and to find their multiplicities.

Clearly there is only one singularity of type I. There are  $(\ell - 1)$  points of type  $(1, \beta)$  with  $\beta^\ell = 1$  and  $\beta \neq 1$ . So, there are also  $(\ell - 1)$  of type  $(\alpha, 1)$  and  $(\ell - 1)$  of type  $(\alpha, \alpha)$  with  $\alpha^\ell = 1$  and  $\alpha \neq 1$ . In total there are  $3(\ell - 1)$  points of type II.

For points of type III there are  $(\ell - 1)$  choices for  $\alpha \neq 1$ , and thus there are  $(\ell - 2)$  choices for  $\beta$  with  $\beta \neq 1$  and  $\beta \neq \alpha$ . However, not all these choices lead to a valid singular point. We upper bound the number of valid choices in the next lemma.

**Lemma 6** *For every  $\alpha$  with  $\alpha^\ell = 1$  and  $\alpha \neq 1$  there exists a  $\beta$  with  $\beta^\ell = 1$ ,  $\beta \neq \alpha$  and  $\beta \neq 1$  such that  $(\alpha + \beta + 1)^\ell \neq 1$ .*

**Proof:** Suppose the statement is false, and fix an  $\alpha \neq 1$  such that for all  $\beta$  with  $\beta^\ell = 1$  we also have  $(\alpha + \beta + 1)^\ell = 1$ . Let  $H$  be  $\{a \mid a^\ell = 1\}$ , the set of  $\ell$ -th roots of unity. Consider the map,

$$\phi : H \rightarrow H, \phi(\beta) = \alpha + \beta + 1.$$



The key point is that this map has no fixed points. For, if  $\phi(\beta) = \beta$ , then  $\alpha = 1$ , which is not true by assumption. Thus  $\phi$  is a permutation of  $H$  which is a product of transpositions of the form  $(\beta, 1 + \alpha + \beta)$ . Therefore  $\phi$  must permute an even number of points, which contradicts the fact that  $\ell$  is odd.  $\square$

From this lemma it follows that, given  $\alpha$ , there are at most  $(\ell - 3)$  possible choices for  $\beta$ . We can not guarantee that each of these is valid, so we can only upper bound the points of type III by  $\leq (\ell - 1)(\ell - 3)$ . There are cases when this bound is tight.

The next Lemma helps us determine when  $m_P(f_t)$  is equal to  $2^i$  and when it is  $2^i + 1$ .

**Lemma 7** *Let  $P = (\alpha, \beta)$  be a singular point of  $f_t$ , then  $F_{2^i} = 0$  if and only if one of the following holds.*

- (1)  $P$  is of Type I
- (2)  $P$  is of Type II.A
- (3)  $P$  is of Type III.A
- (4)  $P$  is of Type III.B and  $\alpha/\beta$  and  $\beta/\lambda \in GF(2^i)$ . In this case, we have  $1 < \gcd(\ell, 2^i - 1) < \ell$ .

**Proof:** We have to check when  $\alpha^{t-2^i} + \lambda^{t-2^i} = 0$ . Substituting  $t = 2^i\ell + 1$  in the formula we get  $\alpha^{1-2^i} = \lambda^{1-2^i}$ , or  $\alpha^{2^i-1} = \lambda^{2^i-1}$ . Now reasoning with  $\beta^{t-2^i} + \lambda^{t-2^i} = 0$  we also obtain that either  $\beta^{2^i-1} = \lambda^{2^i-1}$ . So  $F_{2^i} = 0$  if and only if  $\alpha^{2^i-1} = \beta^{2^i-1} = \lambda^{2^i-1}$ . Consequently,  $F_{2^i} = 0$  if and only if  $(\alpha/\beta)^{2^i-1} = (\beta/\lambda)^{2^i-1}$ .

If  $P$  is of Type I or II.A or III.A, then in fact  $\alpha^{2^i-1} = \beta^{2^i-1} = \lambda^{2^i-1} = 1$ . If  $P$  is of Type II.B then  $F_{2^i} \neq 0$  because certainly one coefficient does not vanish. Finally, suppose  $P$  is of Type III.B, and then we may deduce that  $\alpha = C\beta$  and  $\beta = D\lambda$  for some  $C, D \in GF(2^i)$ . Raising to the  $\ell$ -th power yields that  $C, D$  are  $\ell$ -th roots of unity. Letting  $d = \gcd(\ell, 2^i - 1)$ , then  $C, D$  are  $d$ -th roots of unity. Because  $C$  and  $D$  cannot be 1, we must have  $d > 1$ . If  $d = \ell$  then all  $\ell$ -th roots of unity are in  $GF(2^i)$ . Because  $P$  is of Type III.B, at least one of  $\alpha, \beta$  is not in  $GF(2^i)$ , so  $d < \ell$ .  $\square$

Note that if  $\ell = 2^i - 1$  then  $t = 2^i\ell + 1 = 4^i - 2^i + 1$ , which is an exceptional value.

We now list the classification in a table. We let  $w(x, y) = (x+1)(y+1)(x+y)$  so that  $f_t = wg_t$  and  $m_P(f_t) = m_P(g_t) + m_P(w)$ . The values of  $m_P(w)$  are easy to work out for the various singular points  $P$ . The implications of Lemma 7

can be summarized in the following tables.

$$\gcd(\ell, 2^i - 1) = 1$$

Type	Number of Points	$m_P(f_t)$	$m_P(g_t)$
I	1	$2^i + 1$	$2^i - 2$
II	$3(\ell - 1)$	$2^i$	$2^i - 1$
III	$\leq (\ell - 1)(\ell - 3)$	$2^i$	$2^i$

In this case, the Type II points are all of Type II.B, and the Type III points are all of Type III.B.

$$\gcd(\ell, 2^i - 1) = \ell$$

Type	Number of Points	$m_P(f_t)$	$m_P(g_t)$
I	1	$2^i + 1$	$2^i - 2$
II	$3(\ell - 1)$	$2^i + 1$	$2^i$
III	$\leq (\ell - 1)(\ell - 3)$	$2^i + 1$	$2^i + 1$

In this case, the Type II points are all of Type II.A, and the Type III points are all of Type III.A.

The case  $1 < \gcd(\ell, 2^i - 1) < \ell$  is a mixture of the previous two cases because  $f_t(x, y)$  has points with multiplicity  $2^i$  and points with multiplicity  $2^i + 1$ . Nevertheless the upper bounds on the *number* of points still hold.

### 3.2 Bezout's Theorem

One of the central results in our work uses Bezout's theorem, which is a classical result in algebraic geometry and appears frequently in the literature [3].

**Bezout's Theorem:** Let  $r$  and  $s$  be two projective plane curves over a field  $k$  of degrees  $D_1$  and  $D_2$  respectively having no components in common. Then,

$$\sum_P I(P, r, s) = D_1 D_2. \quad (3)$$

The sum runs over all the points  $P = (\alpha, \beta) \in \bar{k} \times \bar{k}$ , and by  $I(P, r, s)$  we understand the intersection multiplicity of the curves  $r$  and  $s$  at the point  $P$ . Notice that if  $r$  or  $s$  does not go through  $P$ , then  $I(P, r, s) = 0$ . Therefore, the sum in (3) runs over the singular points of the product  $rs$ .

Using properties  $I(P, r_1 r_2, s) = I(P, r_1, s) + I(P, r_2, s)$  and  $\deg(r_1 r_2) = \deg(r_1) +$

$\deg(r_2)$  one can generalize Bezout's Theorem to several curves  $f_1, f_2, \dots, f_r$ :

$$\sum_P \sum_{1 \leq i < j \leq r} I(P, f_i, f_j) = \sum_{1 \leq i < j \leq r} \deg(f_i) \deg(f_j). \quad (4)$$

The following property of the intersection multiplicity will be useful for us. It is part of the definition of intersection multiplicity in [3]. We state it as a Corollary.

**Corollary 8**

$$I(P, r, s) \geq m_P(r)m_P(s), \quad (5)$$

*and equality holds if and only if the tangent cones of  $r$  and  $s$  do not share any linear factor.*

Janwa-McGuire-Wilson [5] have computed the intersection multiplicity at points of type II.B assuming the curve  $g_t(x, y)$  factors:

**Lemma 9** *If  $P$  is a point of type II.B and  $g_t(x, y) = r(x, y)s(x, y)$  then  $I(P, r, s) = 0$ .*

#### 4 The Case $\gcd(\ell, 2^i - 1) < \ell$

Jedlicka has proved in [7] (see Theorem 1) that  $g_t(x, y)$  has an absolutely irreducible factor over  $\mathbb{F}_2$  whenever  $\gcd(\ell, 2^i - 1) < \ell$  and  $t$  is not a Gold or Kasami-Welch number. Therefore, if we prove the same for the case  $\gcd(\ell, 2^i - 1) = \ell$  then conjecture 3' is completely proved. This is what we do in the next section.

#### 5 Main Result: Case $\gcd(\ell, 2^i - 1) = \ell$

The principal starting observation when  $\gcd(\ell, 2^i - 1) = \ell$  is to notice that all  $\ell$ -th roots of unity lie in  $GF(2^i)$ . Therefore, all Type II singularities have Type II.A, and all Type III singularities have Type III.A. From Lemma 7, or the table following it, we have that  $F_{2^i} = 0$  at all singular points.

##### 5.1 Preliminary Lemmata

One of the main ideas involved in the proof in this section is that if  $g_t(x, y)$  is irreducible over  $\mathbb{F}_2$  and splits in several factors (over an extension field),

then all factors have the same degree. The next lemma concerns this sort of phenomenon, and its proof can be found in [9] (although it is surely older).

**Lemma 10** *Suppose that  $p(\underline{x}) \in \mathbb{F}_q[x_1, \dots, x_n]$  is of degree  $t$  and is irreducible in  $\mathbb{F}_q[x_1, \dots, x_n]$ . Then there exists  $r \mid t$  and an absolutely irreducible polynomial  $h(\underline{x}) \in \mathbb{F}_{q^r}[x_1, \dots, x_n]$  of degree  $\frac{t}{r}$  such that*

$$p(\underline{x}) = c \prod_{\sigma \in G} \sigma(h(\underline{x})),$$

where  $G = \text{Gal}(\mathbb{F}_{q^r}/\mathbb{F}_q)$  and  $c \in \mathbb{F}_q$ . Furthermore if  $p(\underline{x})$  is homogeneous, then so is  $h(\underline{x})$ .

Here are some more lemmata we will use.

**Lemma 11** *Given  $N \in \mathbb{N}$  the values  $x_1, \dots, x_n$  that maximize the function  $H(x_1, \dots, x_n) = \sum_{\substack{1 \leq i < j \leq n \\ i \neq j}} x_i x_j$  subject to the constraint  $x_1 + \dots + x_n = N$  are  $x_1 = \dots = x_n = N/n$ .*

One more technical result is recorded now, whose proof is trivial.

**Lemma 12** *If  $i > 2$  and  $\ell \mid 2^i - 1$  but  $\ell \neq 2^i - 1$  then the following results hold:*

- (1)  $2^{i-1} + 1 - \ell > 2$ .
- (2)  $\frac{\ell-3}{2^{i+1}} < \frac{1}{4}$ .

**Proof:** Since  $\ell \mid 2^i - 1$  but  $\ell \neq 2^i - 1$ , and both numbers are odd, we certainly have that  $\ell < 2^{i-1} - 1$ . Then  $2^{i-1} - 1 - \ell > 0$  so  $2^{i-1} + 1 - \ell > 2$ , thus (1) holds.

For (2) we have that  $\ell < 2^{i-1} - 1 < 2^{i-1} + 3$  which implies  $\frac{\ell-3}{2^{i-1}} < 1$  which certainly implies  $\frac{\ell-3}{2^{i+1}} < \frac{1}{4}$ .  $\square$

## 5.2 A Warm-Up Case

**Theorem 13** *Suppose that  $g_t(x, y)$  is irreducible over  $\mathbb{F}_2$  and  $\ell \mid 2^i - 1$  but  $\ell \neq 2^i - 1$ . Then  $g_t(x, y)$  can not split in two factors  $g_1$  and  $g_2$  with  $\deg(g_1) = \deg(g_2)$ .*

**Proof:** We apply Bezout's Theorem, which states

$$\sum_{P \in \text{Sing}(g_t)} I(P, g_1, g_2) = \deg(g_1) \deg(g_2).$$

By Lemma 7 we know that  $F_{2^i} = 0$ . Since the tangent cones have different lines by Lemma 4, Corollary 8 tells us that the left hand side is equal to  $\sum_{P \in \text{Sing}(g_t)} m_P(g_1)m_P(g_2)$ . Using the table of singularities described in Section 3 for  $\ell \mid 2^i - 1$  we get

$$\sum_{P \in \text{Sing}(g_t)} m_P(g_1)m_P(g_2) \leq (2^{i-1} - 1)^2 + 3(\ell - 1)2^{2i-2} + (\ell - 1)(\ell - 3)2^{i-1}(2^{i-1} + 1). \quad (6)$$

Since the degrees of both components are the same, the right hand side of Bezout's Theorem is exactly,

$$(2^{i-1}\ell - 1)^2 = 2^{2i-2}\ell^2 - 2\ell 2^{i-1} + 1. \quad (7)$$

Let us compare (7) and (6). If (7) > (6), we have won, and this happens if and only if,

$$2^{2i-2}(-\ell + 1) + 2^{i-1}(\ell^2 - 2\ell + 1) < 0 \quad (8)$$

which is equivalent to

$$2^{i-1}(\ell - 1) > (\ell^2 - 2\ell + 1) = (\ell - 1)^2. \quad (9)$$

So we conclude that the condition for (7) > (6) is

$$2^{i-1} > (\ell - 1) \quad (10)$$

which is true by Lemma 12 part (1).  $\square$

Remark: Notice that this proof fails when  $\ell = 2^i - 1$ , as it should.

The key idea in the previous proof is to compare (7) and (6). In the next result we have a sharper bound which will be very useful for further results.

**Lemma 14** *If  $\ell \mid 2^i - 1$  but  $\ell \neq 2^i - 1$ , then*

$$\deg(g_t)^2 > \sum_{P \in \text{Sing}(g_t)} m_P(g_t)^2.$$

**Proof:** Suppose not. Then,

$$\begin{aligned} \deg(g_t)^2 &= (2^i\ell - 2)^2 \\ &\leq \sum_{P \in \text{Sing}(g_t)} m_P(g_t)^2 \\ &\leq (2^i - 2)^2 + (3\ell - 3)2^{2i} + (\ell - 1)(\ell - 3)(2^i + 1)^2 \end{aligned}$$

where the last inequality is obtaining using the table of singularities described in section 3. After rearrangement we obtain,

$$0 \leq 2^{2i} + \ell^2 2^{i+1} + \ell^2 - \ell 2^{2i} - 4\ell 2^i - 4\ell + 2^{i+1} + 3.$$

Equivalently,

$$0 \leq 2^i(2(\ell - 1)^2 - 2^i(\ell - 1)) + (\ell - 1)(\ell - 3).$$

Dividing by  $(\ell - 1)$  we get

$$2^{i+1}(2^{i-1} - (\ell - 1)) \leq \ell - 3$$

or

$$2^{i-1} - (\ell + 1) \leq \frac{\ell - 3}{2^{i+1}}.$$

However, by Lemma 12 we know that the left hand side is a positive integer and right hand side satisfies  $0 < \frac{\ell-3}{2^{i+1}} \leq 1/4$ , a contradiction.  $\square$

Remark: Again we note that this proof fails if  $\ell = 2^i - 1$ , as it should.

### 5.3 Proof Assuming Irreducibility over $\mathbb{F}_2$

Next we prove Conjecture 3 under the assumption in the title.

**Theorem 15** *If  $g_t(x, y)$  is irreducible over  $\mathbb{F}_2$ , and  $\ell \mid 2^i - 1$  but  $\ell \neq 2^i - 1$ , then  $g_t(x, y)$  is absolutely irreducible.*

**Proof:** Suppose that  $g_t(x, y)$  is irreducible over  $\mathbb{F}_2$ , and that  $g(x, y) = f_1 \cdots f_r$  over some extension field of  $\mathbb{F}_2$ . By Lemma 10, each  $f_i$  has the same degree, which must be  $\deg(g_t)/r$ . If  $r$  is even then by letting  $g_1 = f_1 \cdots f_{r/2}$  and  $g_2 = f_{1+r/2} \cdots f_r$  we are done by Theorem 13. We may therefore assume that  $r$  is odd (although our argument does not use this, and is also valid when  $r$  is even).

We apply (4) obtaining

$$\sum_P \sum_{1 \leq i < j \leq r} I(P, f_j, f_j) = \sum_{1 \leq i < j \leq r} \deg(f_j) \deg(f_j). \quad (11)$$

The sum over  $P$  is over all singular points of  $g_t$ . Since the degree of  $f_i$  is equal to the degree of  $f_j$  then the right hand side is

$$\sum_{1 \leq i < j \leq r} \deg(f_j) \deg(f_j) = \binom{r}{2} \left( \frac{\deg(g_t)}{r} \right)^2 = \frac{r-1}{2r} \deg(g_t)^2. \quad (12)$$

Now we estimate the inner sum on the left hand side of (11). For any  $P \in \text{Sing}(g_t)$ , since  $F_{2^i+1}$  consists of  $2^i + 1$  different lines by Lemma 5, we have  $I(P, f_i, f_j) = m_P(f_i)m_P(f_j)$  for any  $i, j$  by Corollary 8. Therefore

$$\sum_{1 \leq i < j \leq r} I(P, f_j, f_j) = \sum_{1 \leq i < j \leq r} m_P(f_j)m_P(f_j). \quad (13)$$

We maximize (13) using Lemma 11. We obtain the upper bound

$$\sum_{1 \leq i < j \leq r} m_P(f_j)m_P(f_j) \leq \binom{r}{2} \left( \frac{m_P(g_t)}{r} \right)^2 = \frac{r-1}{2r} m_P(g_t)^2. \quad (14)$$

We denote by  $I, II$  and  $III$  the set of singular points of type I, II and III respectively. Then left hand side in (11) is equal to

$$\begin{aligned} & \sum_{P \in I} \sum_{1 \leq i < j \leq r} I(P, f_j, f_j) + \sum_{P \in II} \sum_{1 \leq i < j \leq r} I(P, f_j, f_j) + \\ & \sum_{P \in III} \sum_{1 \leq i < j \leq r} I(P, f_j, f_j) \stackrel{(14)}{\leq} \\ & \sum_{P \in I} \frac{r-1}{2r} m_P(g)^2 + \sum_{P \in II} \frac{r-1}{2r} m_P(g)^2 + \sum_{P \in III} \frac{r-1}{2r} m_P(g)^2 \leq \\ & \frac{r-1}{2r} \left( (2^i - 2)^2 + (2^i)^2(3\ell - 3) + (2^i + 1)^2(\ell - 1)(\ell - 3) \right). \quad (15) \end{aligned}$$

Once again, the last inequality is thanks to the table with classification of singularities given in section 3 for  $\ell \mid 2^i - 1$ . If (12)  $>$  (15) then we have won. After canceling the factors of  $(r-1)/2r$ , the inequality (12)  $>$  (15) is

$$(2^i \ell - 2)^2 > (2^i - 2)^2 + (2^i)^2(3\ell - 3) + (2^i + 1)^2(\ell^2 - 4\ell + 3)$$

which is true because it is exactly the same inequality as that in the proof of Lemma 14.  $\square$

#### 5.4 Proof of Conjecture 3'

In this section we will finally prove Conjecture 3'.

**Theorem 16** *If  $\ell \mid 2^i - 1$  but  $\ell \neq 2^i - 1$ , then  $g_t(x, y)$  always has an absolutely irreducible factor over  $\mathbb{F}_2$ .*

Proof: Suppose  $g_t = f_1 \cdots f_r$  is the factorization into irreducible factors over  $\mathbb{F}_2$ . Let  $f_k = f_{k,1} \cdots f_{k,n_k}$  be the factorization of  $f_k$  into  $n_k$  absolutely irreducible factors. Each  $f_{k,j}$  has degree  $\deg(f_k)/n_k$ , by Lemma 10.

Let us prove an auxiliary result.

**Lemma 17** *All  $\mathbb{F}_2$ -irreducible components  $f_k(x, y)$  of  $g_t(x, y)$  satisfy the following conditions:*

•

$$\deg(f_k)^2 \leq \sum_{P \in \text{Sing}(g_t)} m_P(f_k)^2. \quad (16)$$

•

$$\sum_{1 \leq i < j \leq n_k} m_P(f_{k,i}) m_P(f_{k,j}) \leq m_P(f_k)^2 \frac{n_k - 1}{2n_k}. \quad (17)$$

**Proof:** Applying Bezout's theorem to  $f_k$  gives

$$\sum_{1 \leq i < j \leq n_k} \sum_{P \in \text{Sing}(f_k)} I(P, f_{k,i}, f_{k,j}) = \sum_{1 \leq i < j \leq n_k} \deg(f_{k,i}) \deg(f_{k,j}) = \deg(f_k)^2 \frac{n_k - 1}{2n_k}. \quad (18)$$

Since for every  $i, j \in \{1, \dots, n_k\}$  the tangent cones of  $f_{k,i}$  and  $f_{k,j}$  consist of different lines by Lemma 4, the left hand side of (18) is

$$\sum_{1 \leq i < j \leq n_k} \sum_{P \in \text{Sing}(f_k)} I(P, f_{k,i}, f_{k,j}) = \sum_{P \in \text{Sing}(f_k)} \sum_{1 \leq i < j \leq n_k} m_P(f_{k,i}) m_P(f_{k,j}) \quad (19)$$

by Corollary 8. We fix  $P$  a singular point. Applying Lemma 11 to

$$\sum_{1 \leq i < j \leq n_k} m_P(f_{k,i}) m_P(f_{k,j})$$

subject to  $\sum_{i=1}^{n_k} m_P(f_{k,i}) = m_P(f_k)$  we get that

$$\sum_{1 \leq i < j \leq n_k} m_P(f_{k,i}) m_P(f_{k,j}) \leq m_P(f_k)^2 \frac{n_k - 1}{2n_k}$$

which proves (17). Summing over  $P$  then proves (16).  $\square$

**Proof of Theorem 16:**

We apply Bezout's Theorem (equation (4)) one more time to the product

$$f_1 f_2 \dots f_r = (f_{1,1} \dots f_{1,n_1})(f_{2,1} \dots f_{2,n_2}) \dots (f_{r,1} \dots f_{r,n_r}).$$

The sum of the intersection multiplicities (left hand side of equation (4)) can be written

$$\sum_{k=1}^r \sum_{1 \leq i < j \leq n_k} \sum_{P \in \text{Sing}(g_t)} I(P, f_{k,i}, f_{k,j}) + \sum_{1 \leq k < l \leq r} \sum_{\substack{1 \leq i \leq n_k \\ 1 \leq j \leq n_l}} \sum_{P \in \text{Sing}(g_t)} I(P, f_{k,i}, f_{l,j})$$



where the first term is for factors within each  $f_k$ , and the second term is for cross factors between  $f_k$  and  $f_l$ . Since for every  $k$  and  $i$  the tangent cones of the  $f_{k,i}$  consist of different lines by Lemma 4, the previous sums can be written

$$\sum_{P \in \text{Sing}(g_t)} \left[ \sum_{k=1}^r \sum_{1 \leq i < j \leq n_k} m_P(f_{k,i}) m_P(f_{k,j}) + \sum_{1 \leq k < l \leq r} \sum_{\substack{1 \leq i \leq n_k \\ 1 \leq j \leq n_l}} m_P(f_{k,i}) m_P(f_{l,j}) \right]. \quad (20)$$

Note that

$$\begin{aligned} (m_P(g_t))^2 &= \left( \sum_{k=1}^r m_P(f_k) \right)^2 \\ &= \sum_{k=1}^r m_P(f_k)^2 + 2 \left( \sum_{1 \leq k < l \leq r} m_P(f_k) m_P(f_l) \right) \\ &= \sum_{k=1}^r m_P(f_k)^2 + 2 \sum_{1 \leq k < l \leq r} \left( \sum_{i=1}^{n_k} m_P(f_{k,i}) \right) \left( \sum_{j=1}^{n_l} m_P(f_{l,j}) \right) \\ &= \sum_{k=1}^r m_P(f_k)^2 + 2 \sum_{1 \leq k < l \leq r} \sum_{\substack{1 \leq i \leq n_k \\ 1 \leq j \leq n_l}} m_P(f_{k,i}) m_P(f_{l,j}). \end{aligned}$$

Substituting, (20) becomes

$$\sum_{P \in \text{Sing}(g_t)} \left[ \sum_{k=1}^r \sum_{1 \leq i < j \leq n_k} m_P(f_{k,i}) m_P(f_{k,j}) + \frac{1}{2} \left( m_P(g_t)^2 - \sum_{k=1}^r m_P(f_k)^2 \right) \right]. \quad (21)$$

Substituting (17) this is

$$\leq \sum_{P \in \text{Sing}(g_t)} \left[ \sum_{k=1}^r m_P(f_k)^2 \frac{n_k - 1}{2n_k} + \frac{1}{2} \left( m_P(g_t)^2 - \sum_{k=1}^r m_P(f_k)^2 \right) \right] \quad (22)$$

$$= \frac{1}{2} \sum_{P \in \text{Sing}(g_t)} \left[ m_P(g_t)^2 - \sum_{k=1}^r \frac{m_P(f_k)^2}{n_k} \right]. \quad (23)$$

On the other hand, the right-hand side of Bezout's Theorem (equation (4)) is

$$\sum_{k=1}^r \sum_{1 \leq i < j \leq n_k} \deg(f_{k,i}) \deg(f_{k,j}) + \sum_{1 \leq k < l \leq r} \sum_{\substack{1 \leq i \leq n_k \\ 1 \leq j \leq n_l}} \deg(f_{k,i}) \deg(f_{l,j}). \quad (24)$$

Since each  $f_{k,i}$  has the same degree for all  $i$ , the first term is equal to

$$\sum_{k=1}^r \deg(f_k)^2 \frac{n_k - 1}{2n_k} = \frac{1}{2} \sum_{k=1}^r \deg(f_k)^2 - \frac{1}{2} \sum_{k=1}^r \frac{\deg(f_k)^2}{n_k}.$$

Note that

$$\begin{aligned}
(\deg(g_t))^2 &= \left( \sum_{k=1}^r \deg(f_k) \right)^2 \\
&= \sum_{k=1}^r \deg(f_k)^2 + 2 \left( \sum_{1 \leq k < l \leq r} \deg(f_k) \deg(f_l) \right) \\
&= \sum_{k=1}^r \deg(f_k)^2 + 2 \sum_{1 \leq k < l \leq r} \left( \sum_{i=1}^{n_k} \deg(f_{k,i}) \right) \left( \sum_{j=1}^{n_l} \deg(f_{l,j}) \right) \\
&= \sum_{k=1}^r \deg(f_k)^2 + 2 \sum_{1 \leq k < l \leq r} \sum_{\substack{1 \leq i \leq n_k \\ 1 \leq j \leq n_l}} \deg(f_{k,i}) \deg(f_{l,j}).
\end{aligned}$$

Substituting both of these into (24) shows that (24) is equal to

$$\frac{1}{2} \left( \deg(g_t)^2 - \sum_{k=1}^r \frac{\deg(f_k)^2}{n_k} \right). \quad (25)$$

Comparing (25) and (23), so far we have shown that Bezout's Theorem implies the following inequality:

$$\deg(g_t)^2 - \sum_{k=1}^r \frac{\deg(f_k)^2}{n_k} \leq \sum_{P \in \text{Sing}(g_t)} \left[ m_P(g_t)^2 - \sum_{k=1}^r \frac{m_P(f_k)^2}{n_k} \right].$$

Finally, using (16) and Lemma 14 to compare both sides term by term, this is a contradiction.  $\square$

## 6 A Counterexample

We have found with MAGMA [1] that  $t = 205$  is a counterexample to Conjecture 3. In this case,  $g_t(x, y)$  factors into two factors over  $\mathbb{F}_2$ . One of the factors is

$$\begin{aligned}
&x^{10} + x^9y + x^9 + x^8y^2 + x^8y + x^8 + x^6y^4 + x^6y^3 + x^6y^2 + x^6y + x^6 + x^5y^5 + x^5 + \\
&x^4y^6 + x^4y^4 + x^4y^3 + x^4y^2 + x^4 + x^3y^6 + x^3y^4 + x^3y^3 + x^3y + x^2y^8 + x^2y^6 + x^2y^4 + \\
&x^2y^2 + x^2 + xy^9 + xy^8 + xy^6 + xy^3 + xy + x + y^{10} + y^9 + y^8 + y^6 + y^5 + y^4 + y^2 + y + 1
\end{aligned}$$

and the other factor has many more terms! Note in this case that  $i = 2$  and  $\ell = 51$ , so  $\gcd(\ell, 2^i - 1) = 3$ . In this case we have  $1 < \gcd(\ell, 2^i - 1) < \ell$ .

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